

# Lecture 2

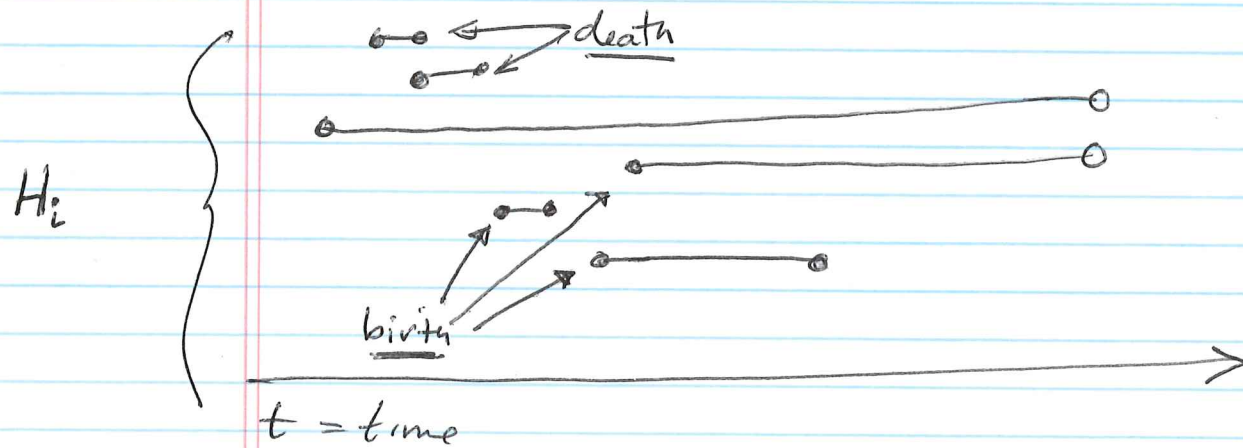
①

## TDA

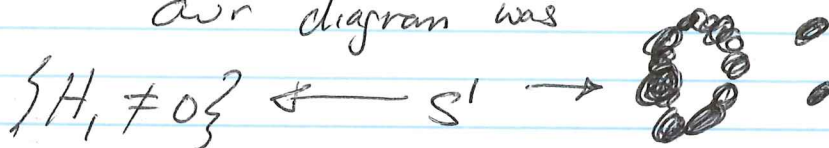
Preamble: For this lecture and next lecture, we will be moving towards the following goal.

"Fundamental Thm of PH" (FTP H)

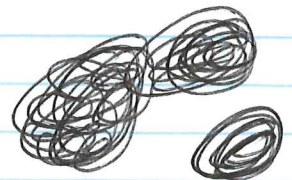
PH is completely encoded by a barcode, consisting of long bars and short bars.



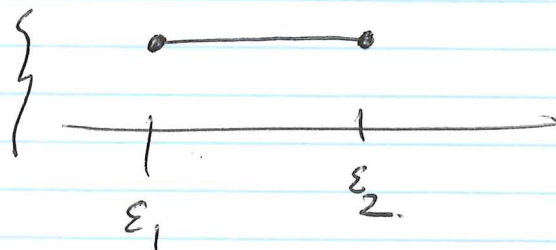
Intuition in our "idea" from last lecture, we saw at some time (birth)  $\epsilon_1$ , our diagram was



and at a later  $\epsilon_2$ , we had



so the  $H_1$  barcode would be



# Lecture 2

①

FTPH is a consequence of the following

[MPID]

Theorem (Artin, 6.13 or Fraleigh, 11.12)

A finitely generated module  $M$  over a Euclidean Domain  $R$  (more generally, over a PID) has a decomposition

$$M \cong R^n \oplus \bigoplus_i R/d_i, \quad \{d_i \neq 0, \text{unit}\}, \quad d_i | d_{i+1}$$

Basic idea: Long bars  $\Leftrightarrow$  free summands  $\Leftrightarrow R^n$   
Short bars  $\Leftrightarrow$  torsion  $\Leftrightarrow \bigoplus R/d_i$

Amazing and Beautiful Fact MPID has very different avatars

① 11.12 (Fraleigh) • A finitely generated Abelian group is  $\cong \mathbb{Z}^n \oplus \bigoplus_i \mathbb{Z}/(p_i^{e_i})$   $p_i$  prime

② 12.7.9 (Artin) • A linear operator on a finite dim vector space can be written (Δ basis) as a block diagonal matrix, with blocks

$$\begin{bmatrix} 0 & & & -a_1 \\ & \ddots & & \vdots \\ & & 1 & 0 \\ & & & \ddots \\ & & & & -a_n \\ & & & & & 1 & -a_n \end{bmatrix} \text{ } \left. \vphantom{\begin{bmatrix} 0 & & & -a_1 \\ & \ddots & & \vdots \\ & & 1 & 0 \\ & & & \ddots \\ & & & & -a_n \\ & & & & & 1 & -a_n \end{bmatrix}} \right\} \text{ square.}$$

A matrix of this type is said to be in rational canonical form: it is the ~~best~~ closest we can get to diagonalize in general  
(not every matrix can be diagonalized)





18 seconds

# Lecture 2

3

## REST OF TODAY'S LECTURE — RATIONAL CANONICAL FORM

### Turbo Lin Alg Review (Next 20 minutes)

field  
↓  
 $k$

Let  $V$  be a vector space of finite dim (vsfd) over  $k$

- DEF •  $\{v_1, \dots, v_k\}$  spans  $V$  if any  $v \in V$  can be written  $v = \sum_{i=1}^k a_i v_i$
- " is linearly indep (LID) if  $\sum_{i=1}^k a_i v_i = 0 \Rightarrow \text{all } a_i = 0$
  - " is a basis if it spans and is LID.

Exercise: Cardinality of a basis (vsfd) is well defined.

DEF A map  $V \xrightarrow{T} V$  is a linear transformation (LT) if  $T(\alpha v_1 + v_2) = \alpha T(v_1) + T(v_2), v_i \in V, \alpha \in k.$

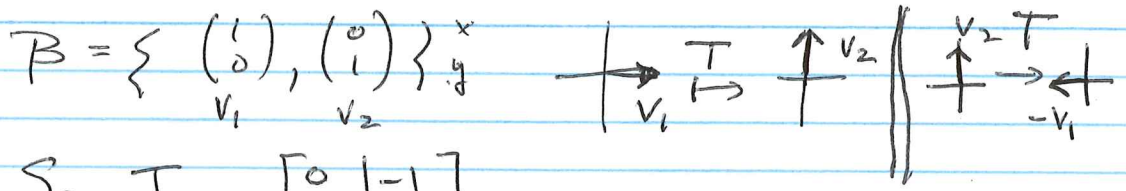
Exercise A LT  $V \rightarrow V \cong k^n$  may be written as an  $n \times n$  matrix, via the procedure

$B = \{v_1, \dots, v_n\}$  basis for  $V$

$$A_B = \left[ \begin{array}{c|c} T(v_1)_B & \dots & T(v_n)_B \end{array} \right]$$

where  $T(v_i)_B$  means apply  $T(v_i)$ , write result in terms of  $B$ .

Example:  $T: \mathbb{R}^2 \rightarrow \mathbb{R}^2$  via  $90^\circ$  rotation.



So  $T_B = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ .

Check that using  $B' = \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ -1 \end{pmatrix} \right\}$  gives

$$T_{B'} = \begin{bmatrix} 1 & -2 \\ 1 & -1 \end{bmatrix}$$



# Lecture 2

(4)

WHO CARES? Suppose  $T$  represents a

$(T)^n = A^n = \begin{bmatrix} \text{very} & \$ \\ \text{expensive} & \end{bmatrix}$  ← Markov process, and we want to do the process  $n$  times

If we choose (or are lucky) a basis where  $A$  is diagonal, it is trivial:  $\begin{pmatrix} x_1 & 0 \\ 0 & x_m \end{pmatrix}^n = \begin{pmatrix} x_1^n & 0 \\ 0 & x_m^n \end{pmatrix}$  😊

DEF  $v$  is an eigenvector for  $T$  if  $(T - \lambda I)v = 0$ ,

$\lambda$  is the eigenvalue, and  $Tv = \lambda v$ , so if  $v$  is part of a basis  $\beta$ ,  $(T(v))_{\beta} = \begin{pmatrix} \lambda \\ 0 \\ \vdots \\ 0 \end{pmatrix}$  (say  $v =$  first basis vect)

Hence, if  $T$  has a basis of eigenvectors,  $T$  is diagonalizable 😊

Exercise Find an invertible  $2 \times 2$  matrix which can't be diagonalized

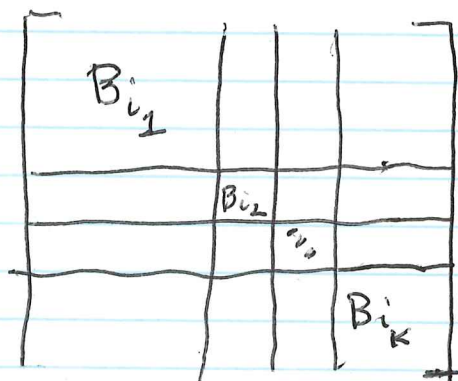
Wake up: Lin alg review is over!

Rational Canonical Form: (Artin §12.7, Lang "Algebra" XIV.2)

DEF: An invariant block is an  $n \times n$  square matrix

$$B_n = \begin{bmatrix} 0 & \dots & 0 & a_1 \\ 1 & & 0 & \vdots \\ 0 & 1 & & \vdots \\ \vdots & & \ddots & \vdots \\ 0 & \dots & 0 & 1 & a_n \end{bmatrix} = I^s \text{ (immediately) below diagonal, } 0 \text{'s elsewhere except last col, } a_i \in k.$$

Theorem Any  $n \times n$  matrix is similar (= after a basis) to a matrix with  $\sum_{j=1}^k i_j = n$  and  $B_{i_j}$  invariant



# Lecture 2

(5)

Pf: Let  $v \in V$ ,  $T$  a linear transform  $V \xrightarrow{T} V$

let  $j$  be the smallest intgr such that

$$T^j(v) \in \text{span} \{v, T(v), T^2(v), \dots, T^{j-1}(v)\} = \beta.$$

Put  $V_{j-1} = \text{span}(\beta)$ ,  $\beta$  is clearly a basis for  $V_{j-1}$

$V_{j-1}$  is a subspace of  $V$ , and writ  $\beta$ ,

$$\left[ T \Big|_{V_{j-1}} \right]_{\beta} = \begin{bmatrix} 0 & \dots & 0 & a_0 \\ 1 & & & \vdots \\ 0 & \ddots & & 0 \\ \vdots & & \ddots & \vdots \\ 0 & \dots & 0 & a_{j-1} \end{bmatrix}, \text{ where } T^j(v) = \sum_{i=0}^{j-1} a_i T^i(v)$$

First col.  $\leftarrow j \times j$  matrix

$v_1 = v \rightarrow T(v) = v_2$

an invariant block.

Now just iterate the process for  $V/V_{j-1}$  □

## Connecting Linear Transforms to $K[t]$ modules:

Thm (Arch, p 476): A linear operator  $T$  on a vstd  $V/k$  and a  $K[t]$  module ( $f$ 's end) are equiv.

Sketch Pf: • Let  $T: V \rightarrow V$  a LT, and  $f(t) = \sum_{i=0}^m a_i t^i$

we make  $V$  into a  $K[t]$ -module via

$$f(t) \cdot v = \sum_{i=0}^m a_i T^i(v).$$

Check this operator satisfies the module requirements.

- If  $V$  is a  $K[t]$  module, then  $K$  acts on  $V$ , so  $V$  is a  $K$ -vector space, and  $\circ t$  maps  $V \rightarrow V$ . The module properties mean  $\circ t$  is a linear op. so define  $T = \circ t$
-



# Lecture 2.

⑥

Thus, we have shown

①  $T$  a linear operator on  $V$  vsfd /  $k$



$V$  a finitely generated  $k[t]$  module

↑  
remember, this is  
a Euclidean Domain!

② Any lin. op  $T$  on a vsfd /  $k$   
has a decomposition into a rational canonical form

$$\begin{bmatrix} B_{i1} & 0 & 0 & 0 \\ 0 & B_{i2} & 0 & 0 \\ 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & B_{ik} \end{bmatrix}$$

With  $B_{ij}$  square  
invariant  
blocks

③ Concrete example: suppose  $V = k[t] / \langle f(t) \rangle$ ,

where  $f(t) = t^n + a_{n-1}t^{n-1} + \dots + a_0$ . Clearly  $V$   
is a vsfd over  $k$ , with basis

$$1, t, \dots, t^{n-1}$$

$$\downarrow$$

$$t \quad t^2 \quad \dots \quad t^n = \sum_{j=0}^{n-1} a_j t^j$$

So matrix is

$$\begin{bmatrix} 0 & 0 & \dots & 0 & -a_0 \\ 1 & 0 & \dots & 0 & \vdots \\ 0 & 1 & \dots & 0 & \vdots \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & 1 & -a_{n-1} \end{bmatrix} \leftarrow \text{invariant!}$$

# Lecture 2

(7)

## TDA

### Jordan Canonical Form

Suppose  $f(t) = (t - \alpha)^n$ ,  $W = K[t] / \langle f(t) \rangle$

We use a slightly different basis

$w_0 :=$  residue class of 1 in  $W$

$$w_i = (t - \alpha)^i w_0$$

$$\text{Then } (t - \alpha)w_0 = w_1$$

$$(t - \alpha)w_{n-2} = w_{n-1} \quad \text{and} \quad (t - \alpha)w_{n-1} = 0$$

$$\text{Hence } (T - \alpha)w_i = w_{i+1} \Rightarrow Tw_i = \alpha w_i + w_{i+1} \quad \left. \vphantom{Tw_i} \right\} i=0, \dots, n-2$$

$$Tw_{n-1} = \alpha w_{n-1}$$

So the matrix for  $T$  is 
$$\begin{bmatrix} \alpha & 0 & \dots & 0 \\ 1 & \alpha & & \\ 0 & 1 & \dots & \\ \vdots & 0 & \ddots & 0 \\ 0 & \vdots & 0 & 1 & \alpha \end{bmatrix} \leftarrow \text{Jordan Block}$$

Notice: If  $K$  is alg closed (say, if  $K = \mathbb{C}$ )

$$\text{then any } f(t) = \prod_{i=1}^n (t - \alpha_i)^{m_i}$$

so we get a representation of  $T$  w/ Jordan blocks

### BONUS: Google's PageRank Algorithm

- (i) Represent the web as a weighted, directed graph  
vertices = website  
edge  $i \rightarrow j$  if site  $i$  has link to site  $j$

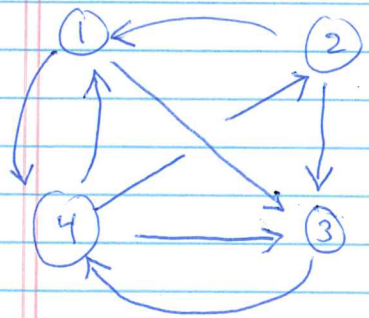
weight: if site  $i$  has  $k$  links = arrows pointing "out"

intuition: If you're at site  $i$ , equally likely to choose one of  $k$  arrows out.  
then each arrow out of  $i$  has weight  $\frac{1}{k}$

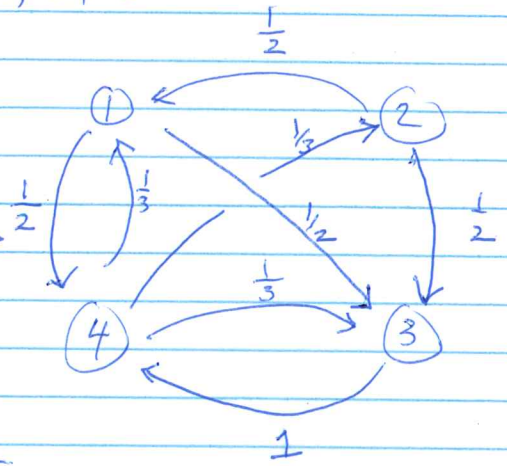


<u>Example</u>	Site 1	links to	3, 4
	2	"	1, 3
<u>Lecture 2</u>	3	"	4
	4	"	1, 2, 3

(8)



weighting,  
we get



(2) From the graph, build a transition matrix:

$$T = \begin{bmatrix} 0 & 1/2 & 0 & 1/3 \\ 0 & 0 & 0 & 1/3 \\ 1/2 & 1/2 & 0 & 1/3 \\ 1/2 & 0 & 1 & 1/3 \end{bmatrix} =: T$$

$T$  is  $n \times n$ ,  $n = \#$  websites  
if  $v_m$  is  $n \times 1$ ,  $\sum v_i = 1$ ,  
then  $T(v) = v_{n+1}$  = probabilities at time  $m+1$

(3) But, there is a probability you don't follow a link, but type in a new URL. Assume jumping to any site is equally likely.

DEF  $G = (1-p)T + p \left( \frac{1}{n} \hat{I} \right)$ ,  $\hat{I}$  is  $n \times n$  matrix with all 1's  
 $p \in [0, 1]$

Rmk:  $G$  is positive + column stochastic  
(all cols sum to 1)

THM (Perron-Frobenius) A positive, column stochastic matrix has 1 as an eval; corr. vect  $v^*$  is all +.

THM If  $v^*$  is as above, and  $v(0) = \frac{1}{n} \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$ , then

$$\lim_{n \rightarrow \infty} G^n (v(0)) = v^* \leftarrow \text{the PageRank!}$$

ALG: To find PageRank, solve  $Gv = 1v$ , scale  $v$  so entries sum to 1.

Problem:  $n$  is in the billions!

[Aside in practice, p. 15]

## Lecture 2

(9)

### Exercises

1. Prove if  $V$  vsfd /  $\mathbb{K}$  then the cardinality of a basis for  $V$  is well defined. DEF  $|\text{Basis}| = \dim V$

2. Find the change of basis matrix  $\Delta_{B'B}: B \rightarrow B'$  for example on  $P_3$ . (notes)

Show  $\Delta_{B'B} \cdot \Delta_{B'B'} = I_2 = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$

3. Find the rational canonical form for  $V = \mathbb{K}[t]/t^3 - 4t^2 + 5t - 2 = f(t)$ . Hint: factor  $f(t)$

4. Diagonalize the matrix  $T$  on  $P_8$  (notes)